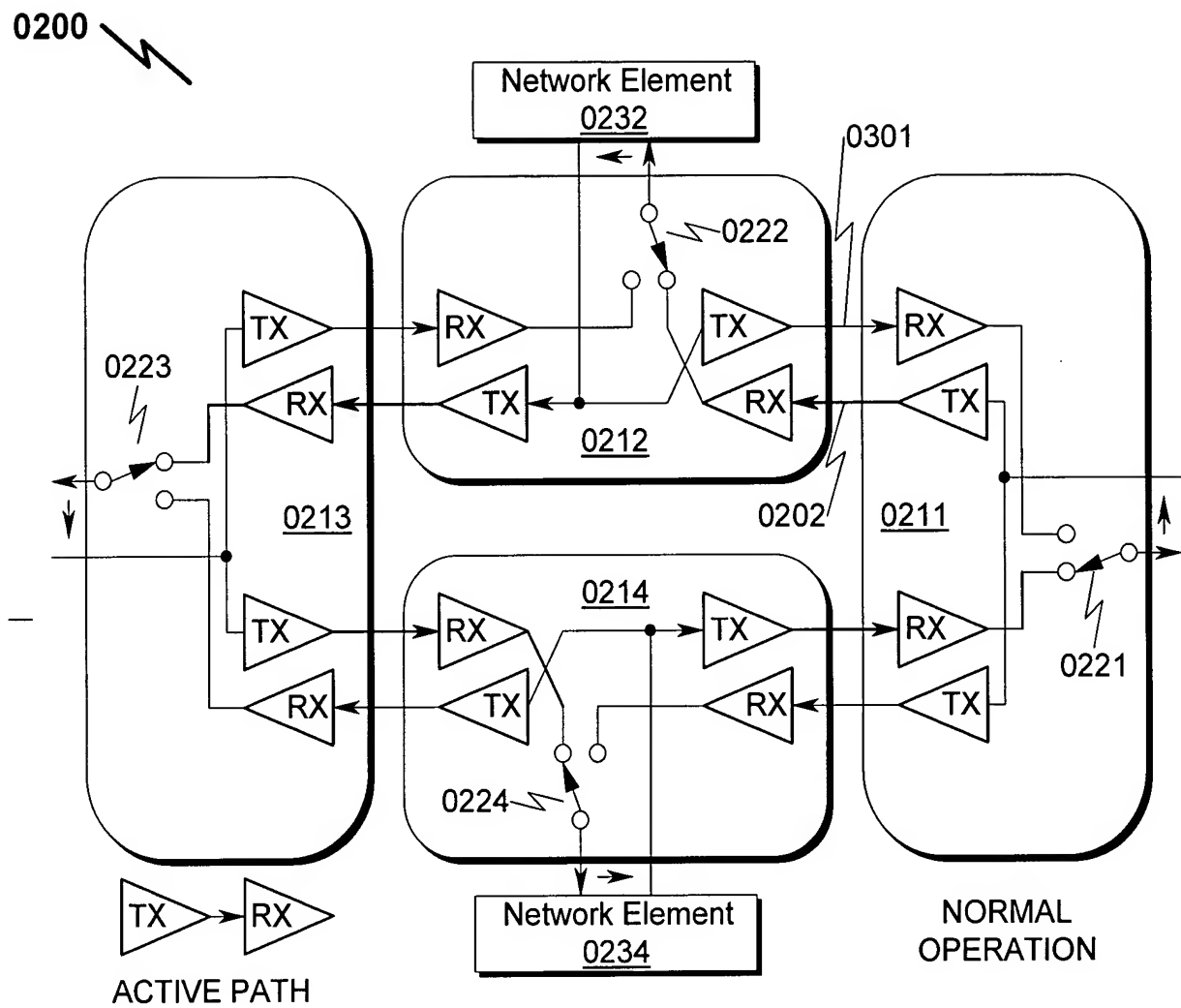


Figure 1



PRIOR ART

Figure 2

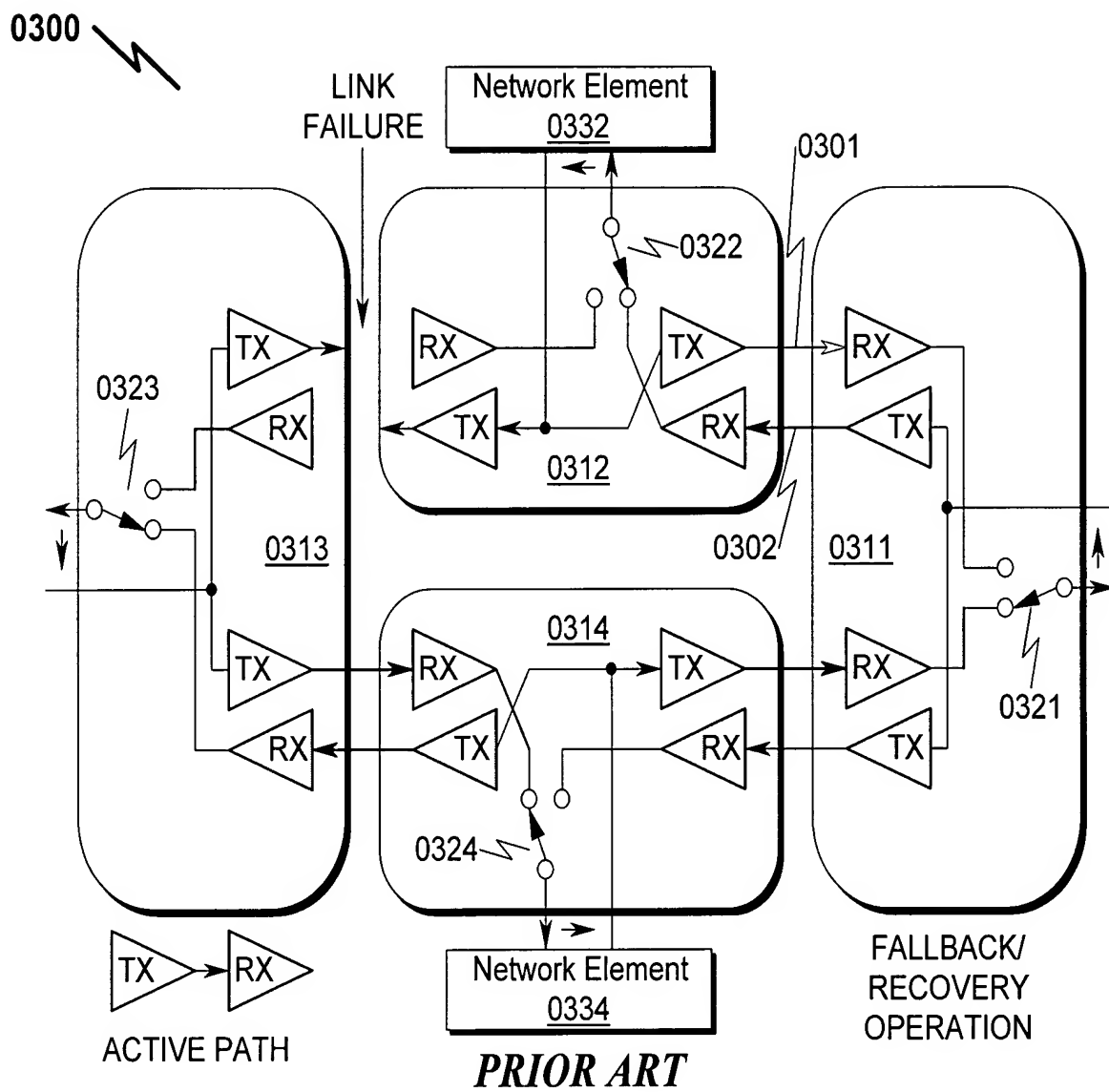
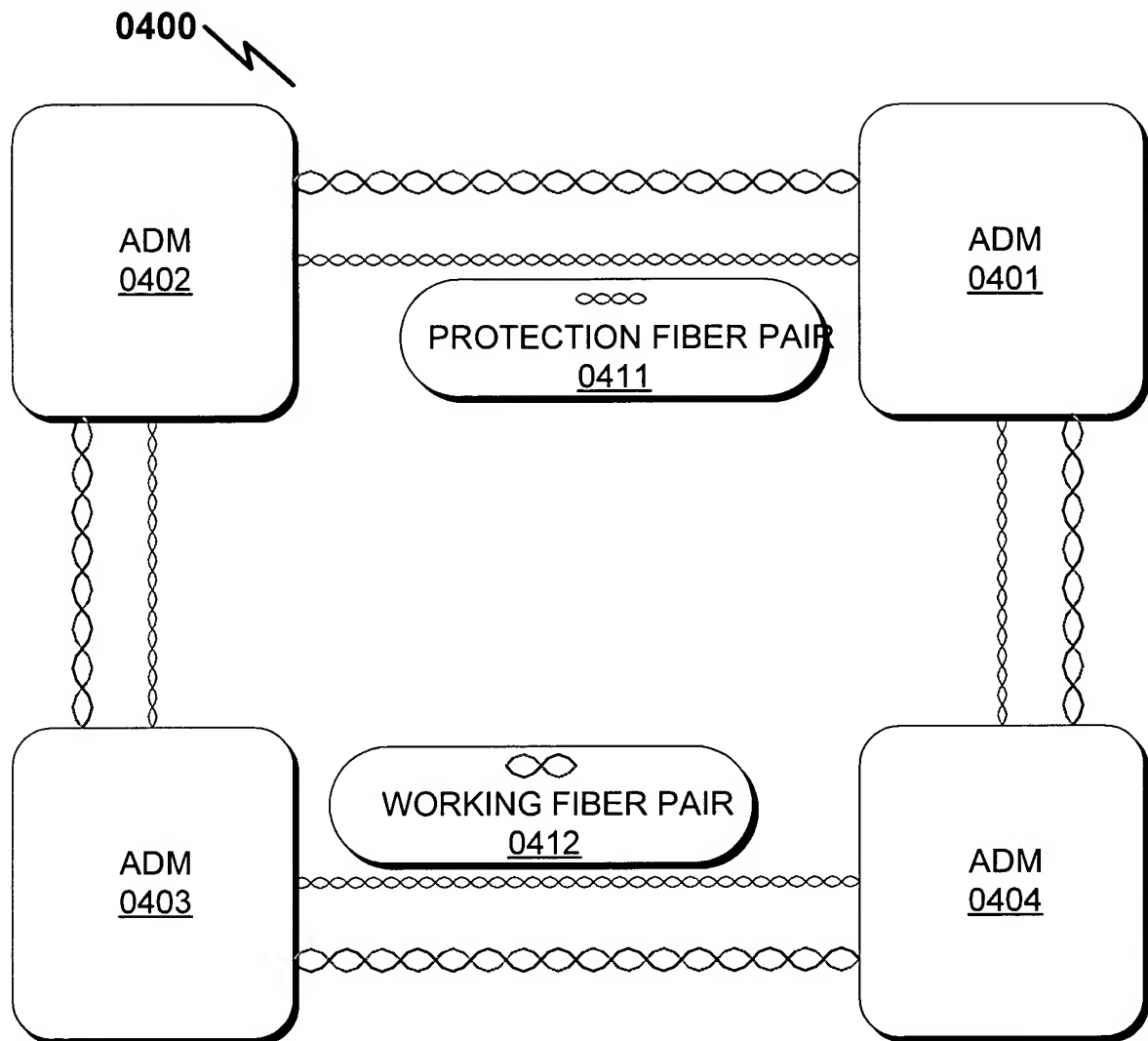


Figure 3



PRIOR ART

Figure 4

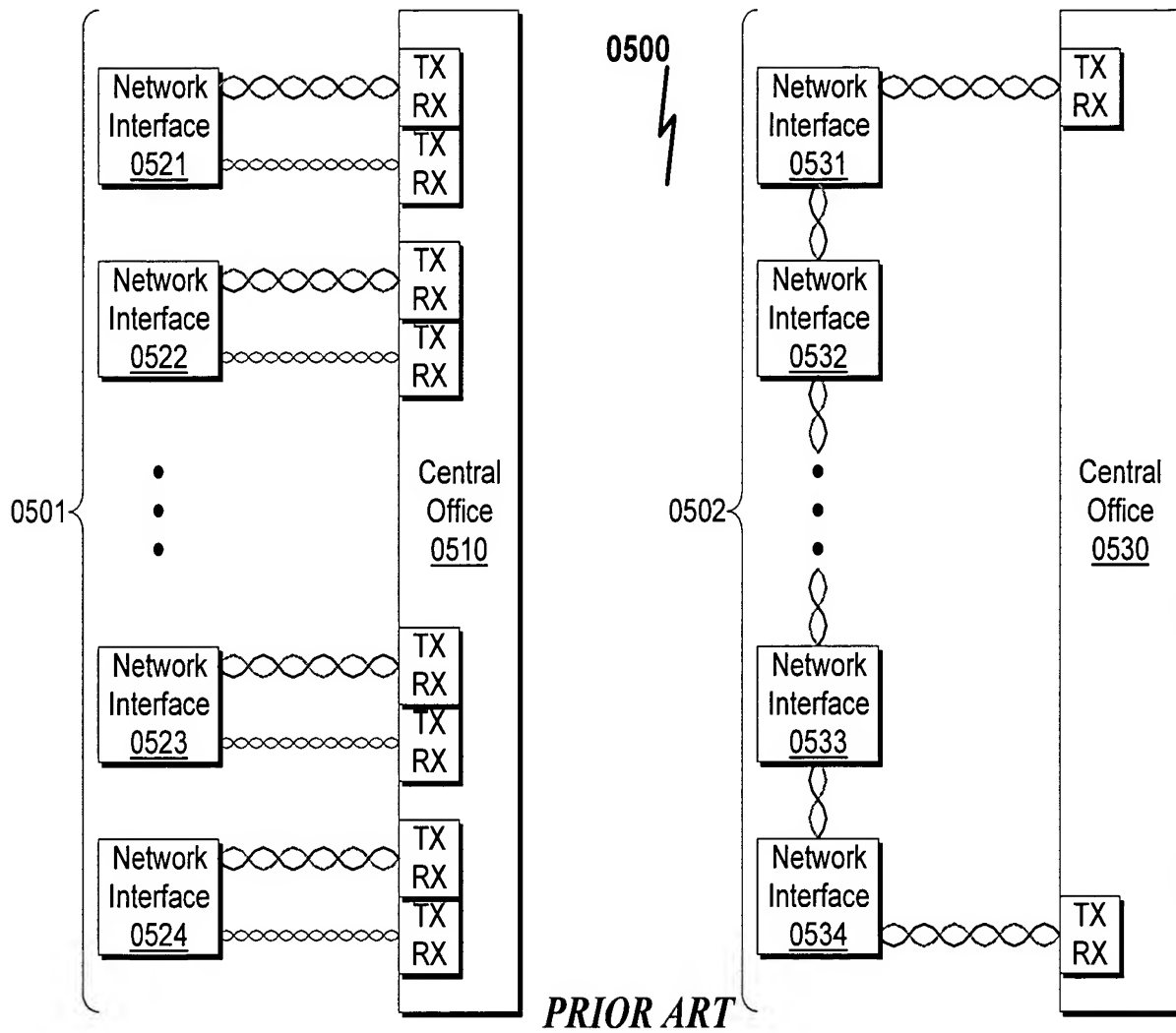


Figure 5

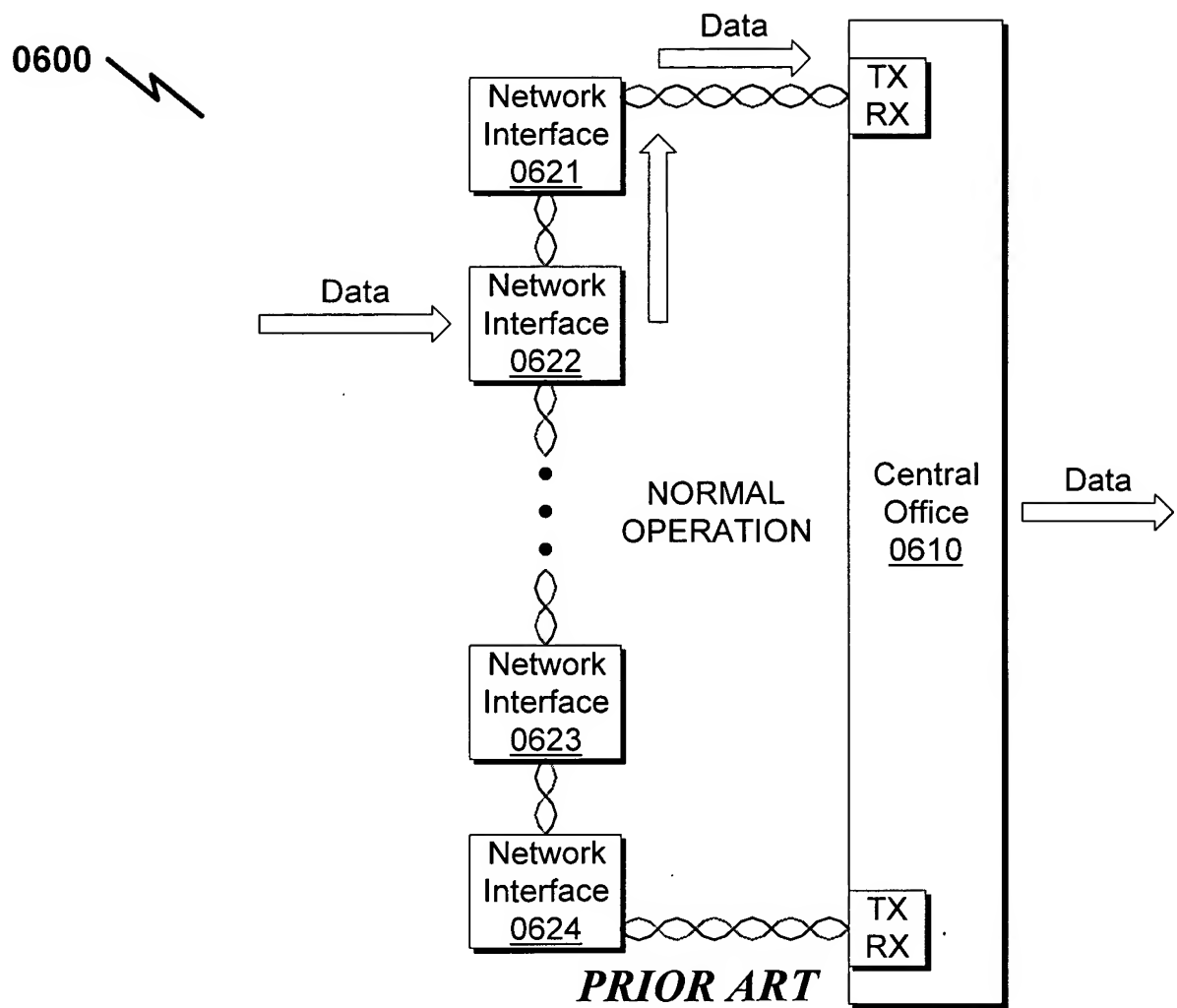


Figure 6

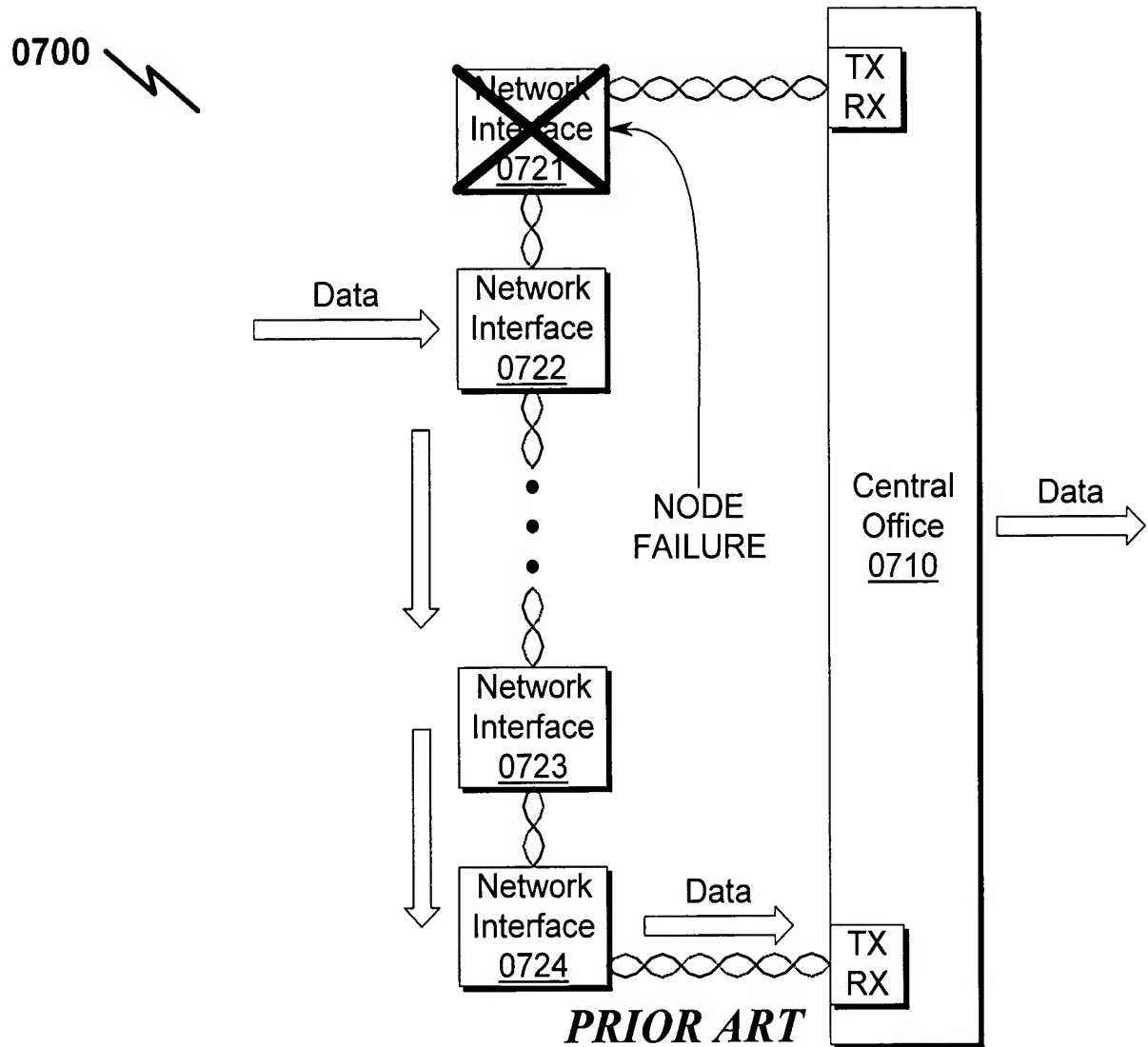


Figure 7

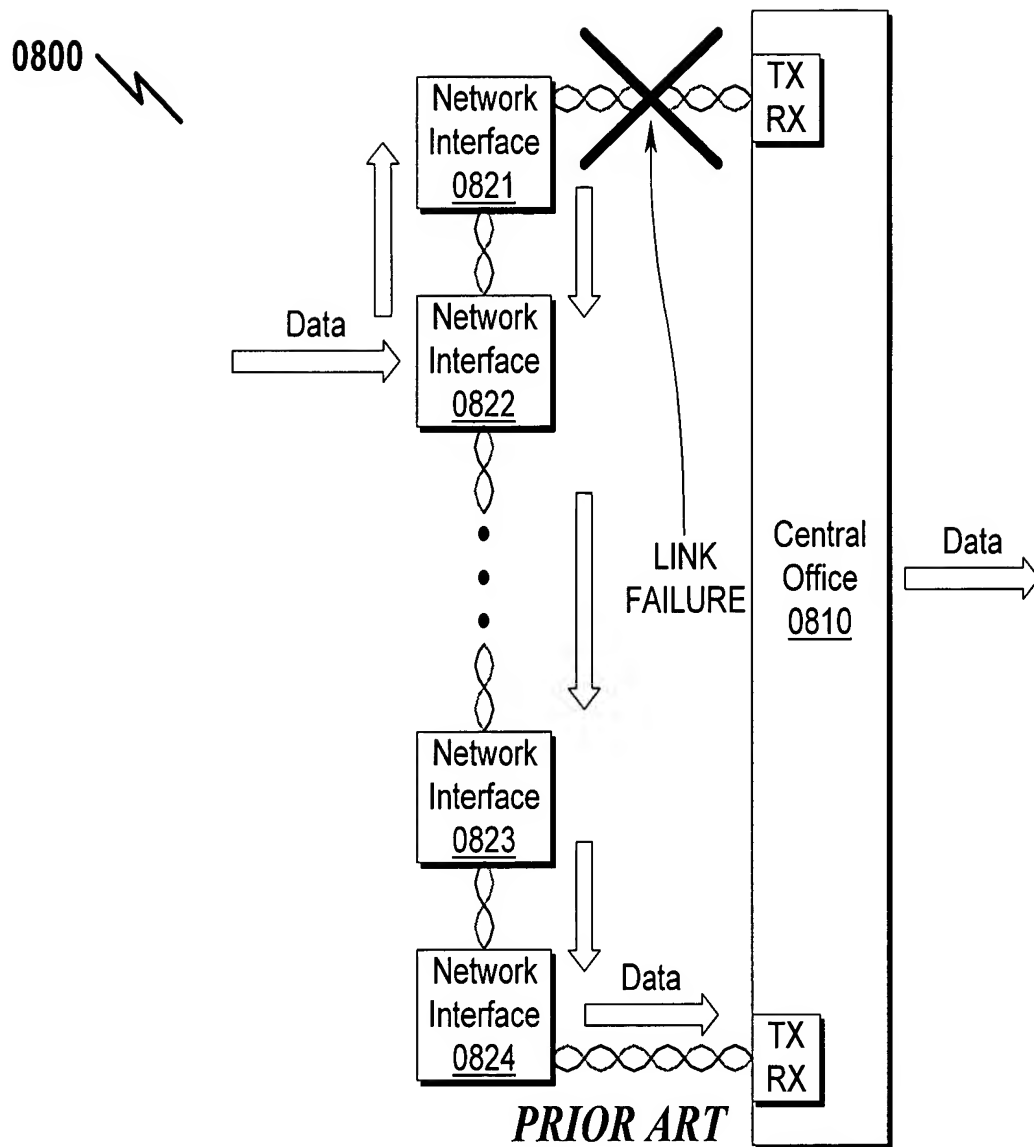


Figure 8

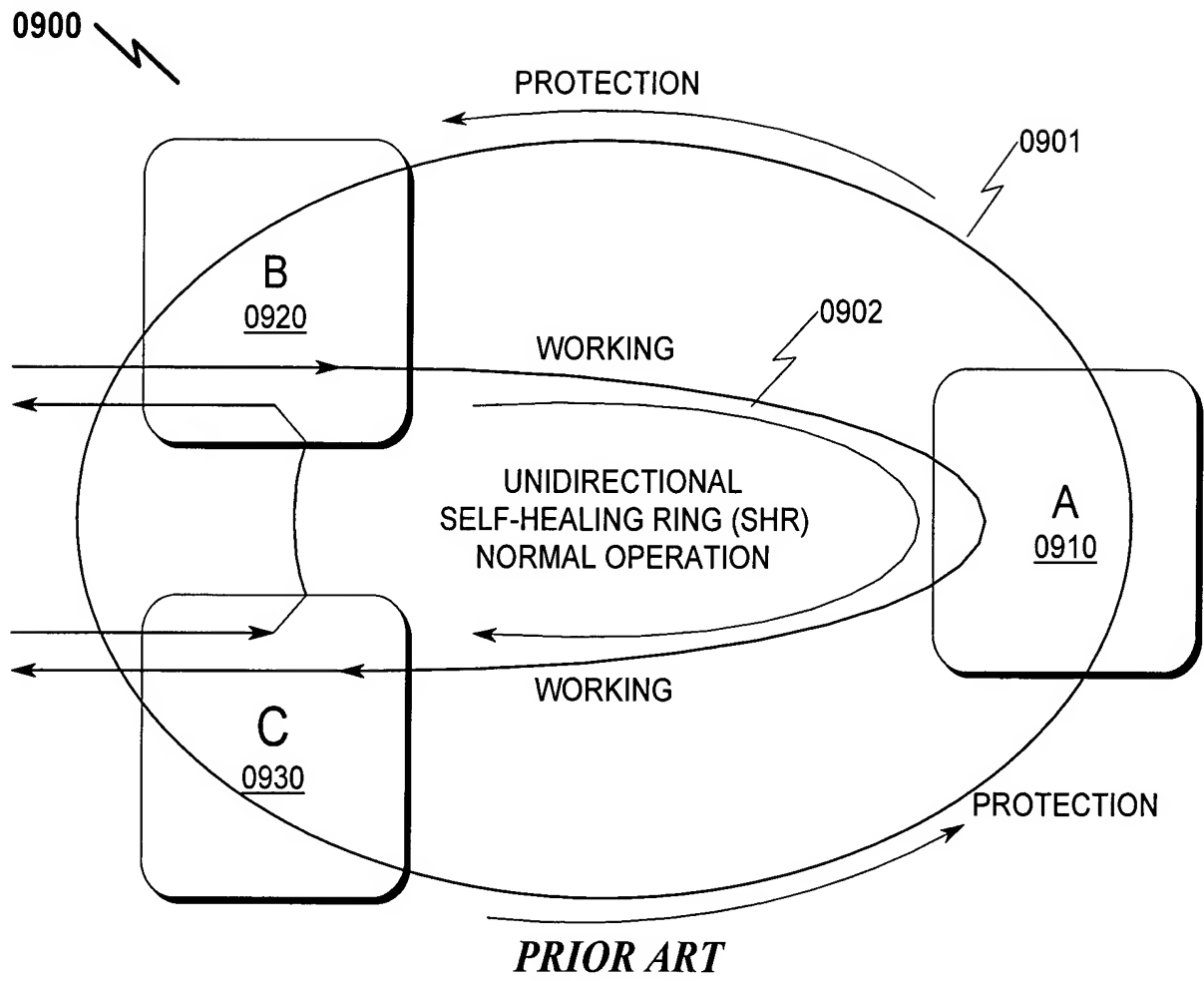


Figure 9

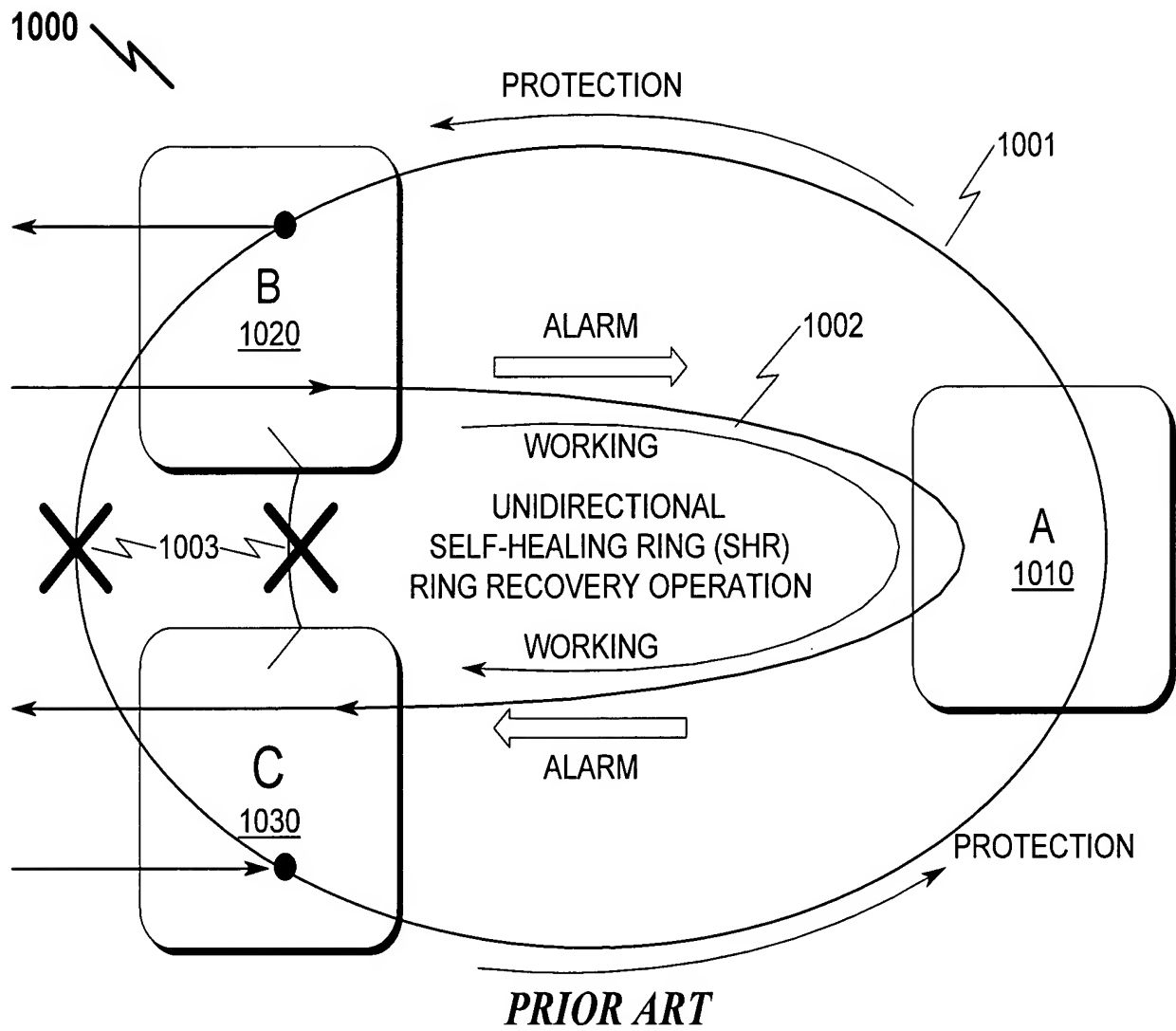


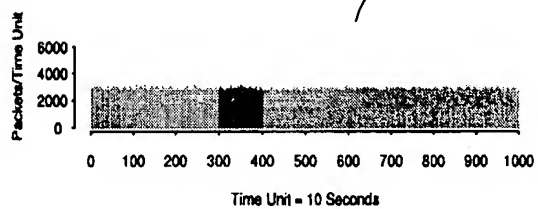
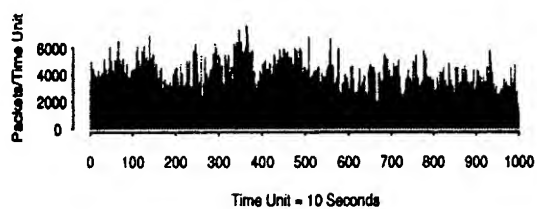
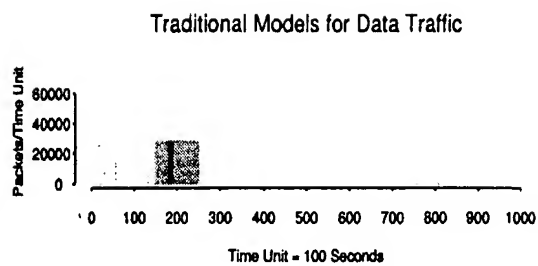
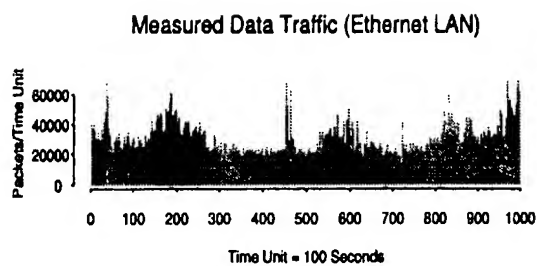
Figure 10

1104

1100

2

1102



1112

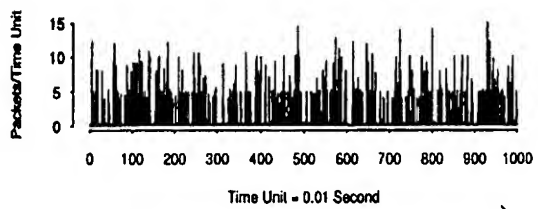
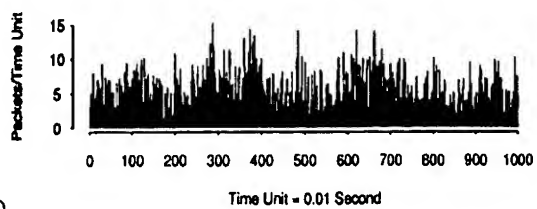
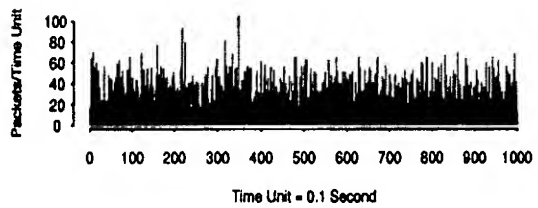
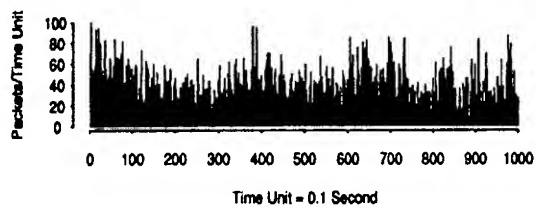
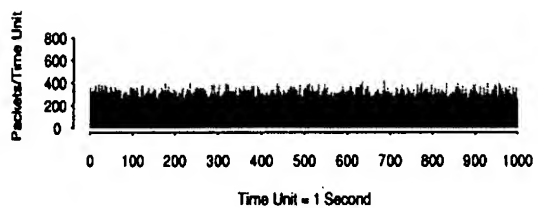
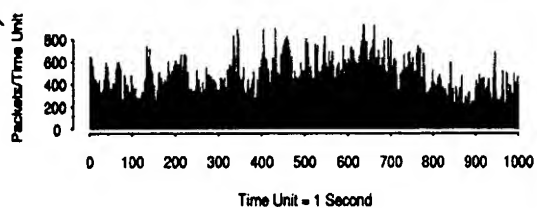


Figure 11

PRIOR ART

1110

1200

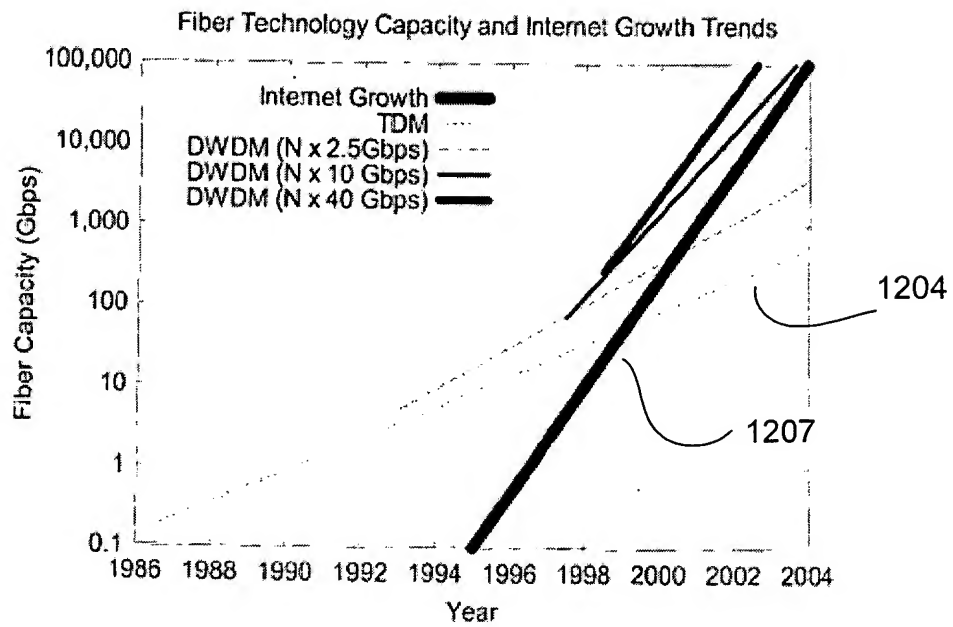


Figure 12

1300

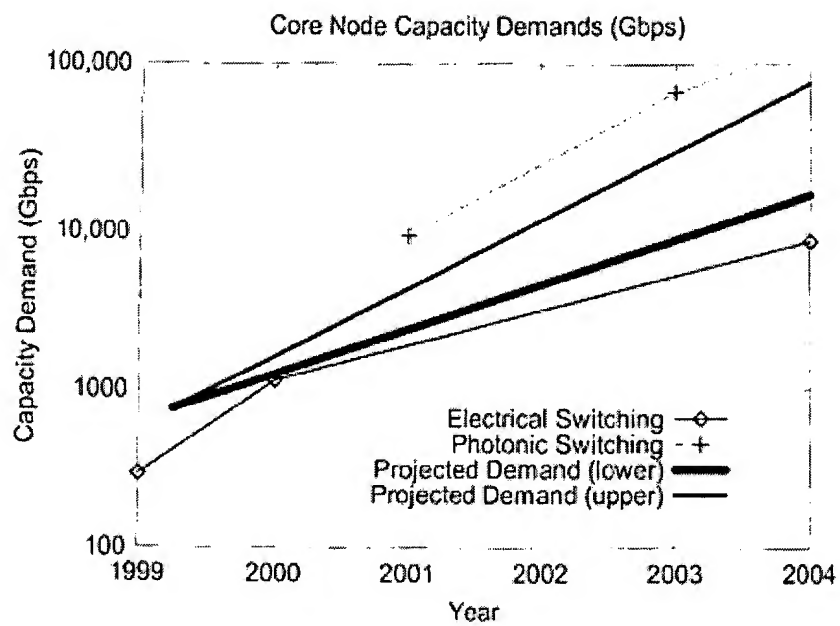


Figure 13

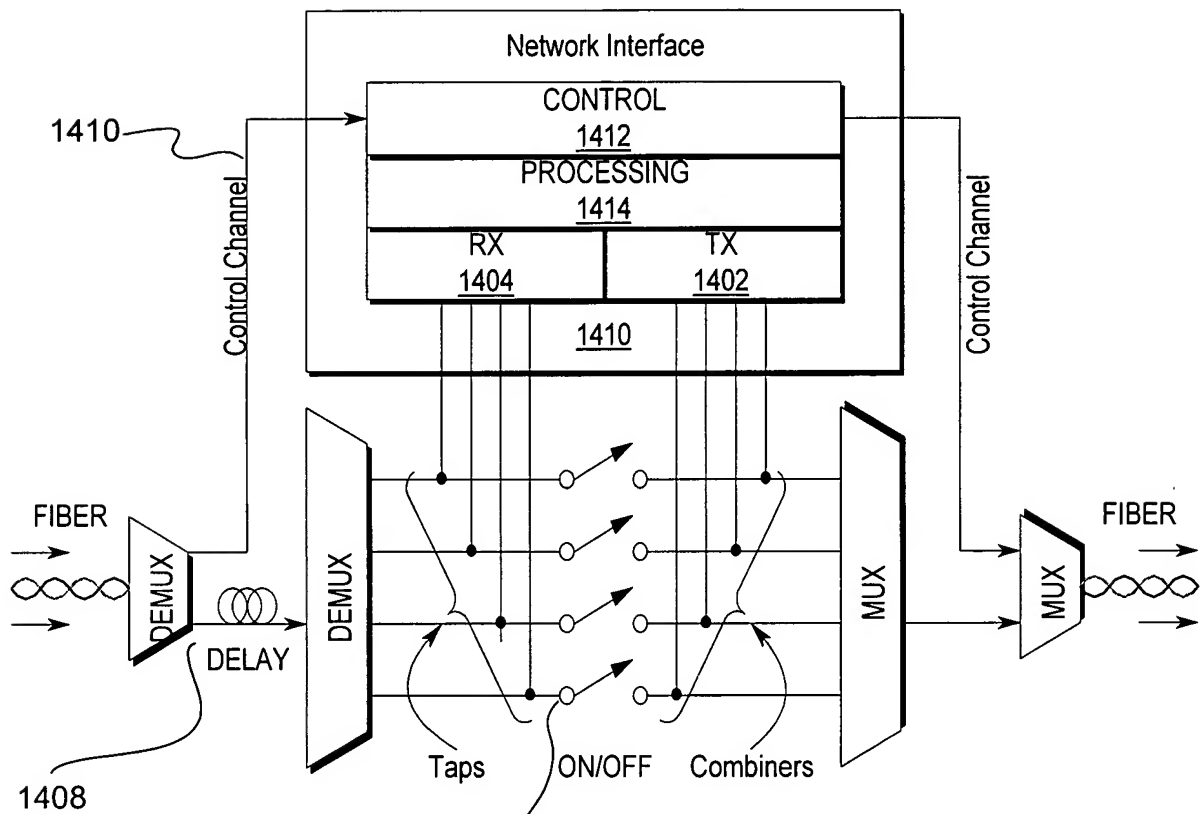


Figure 14
PRIOR ART



Figure 15

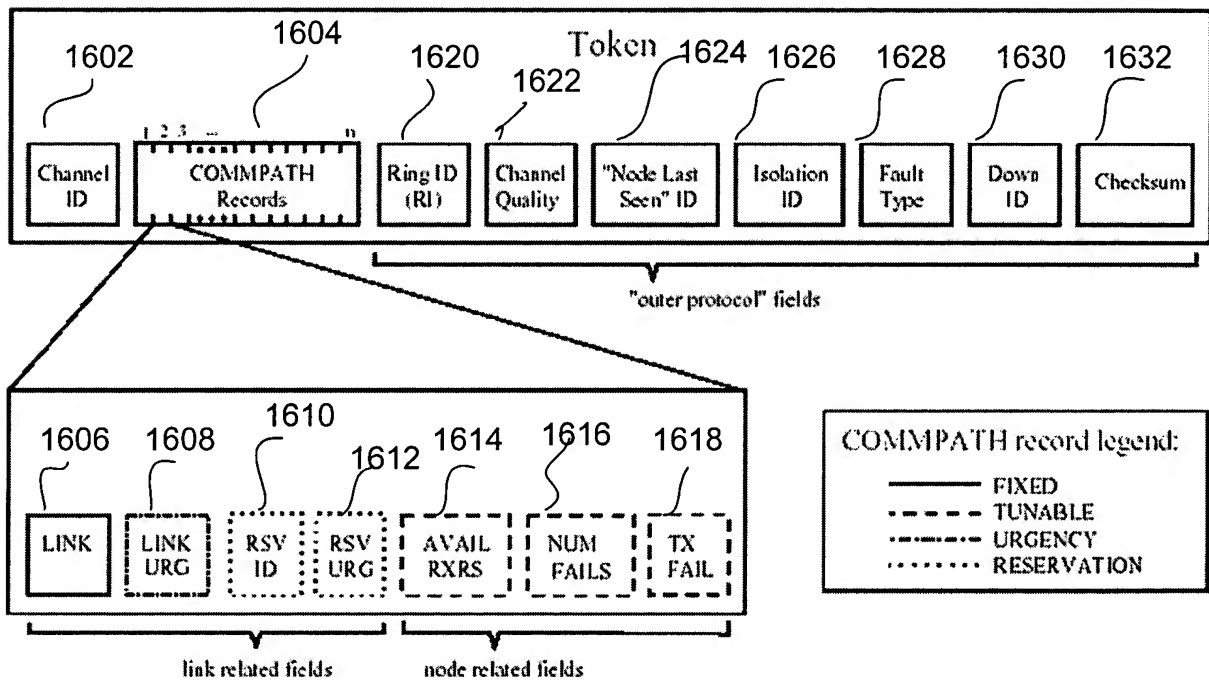
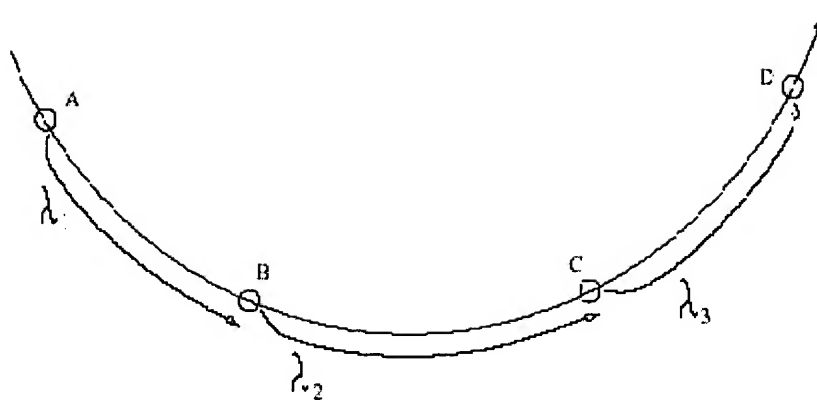
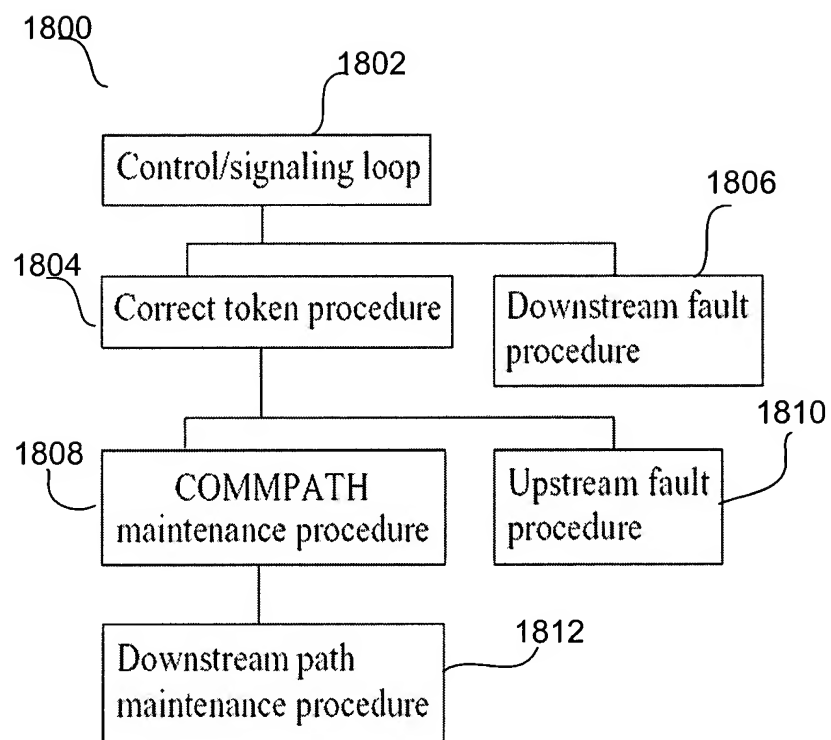


Figure 16



Nodes A, B, and C contend for D's lone RXR

Figure 17



Reference Network protocol procedure dependencies

Figure 18

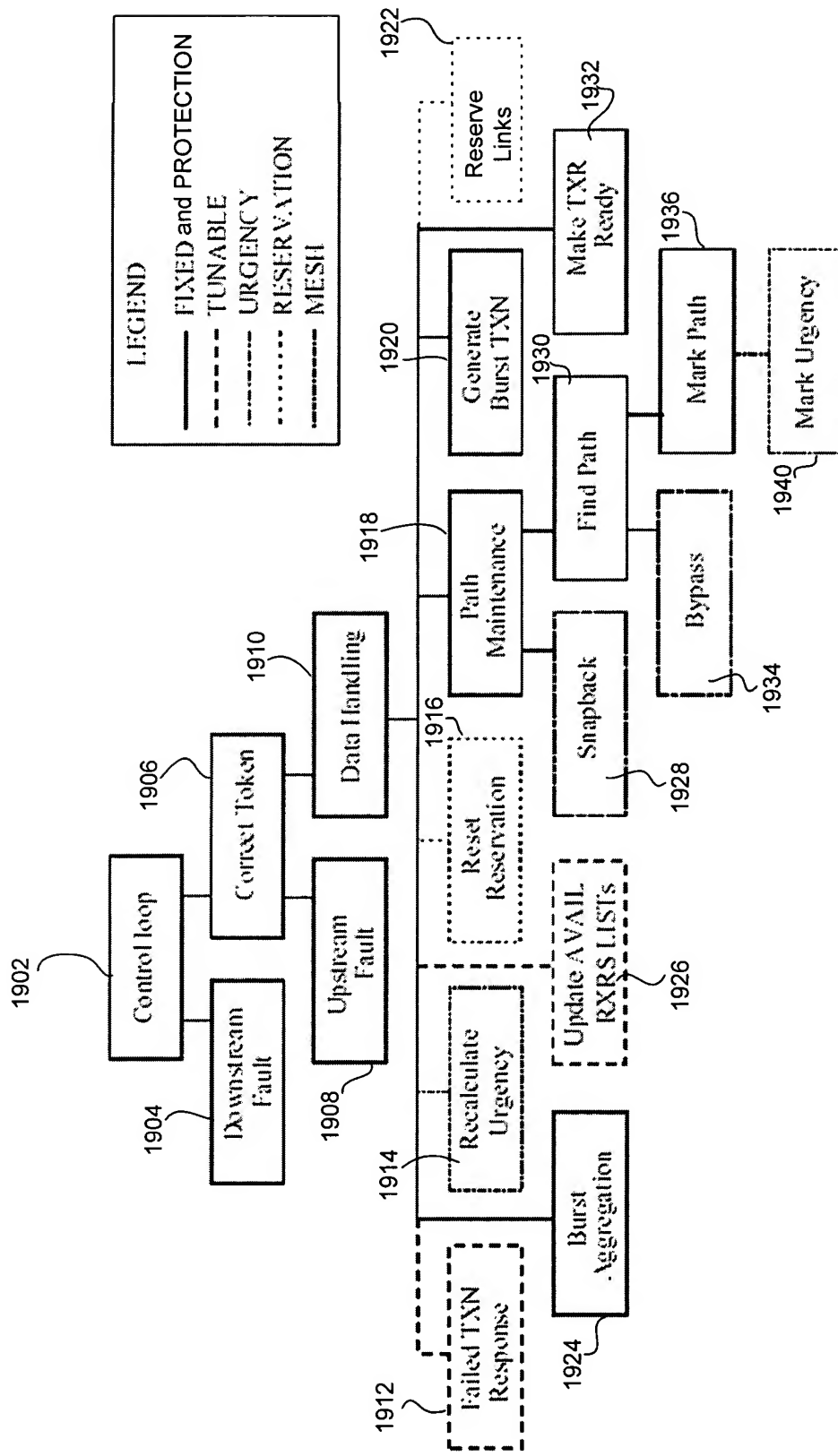


Figure 19

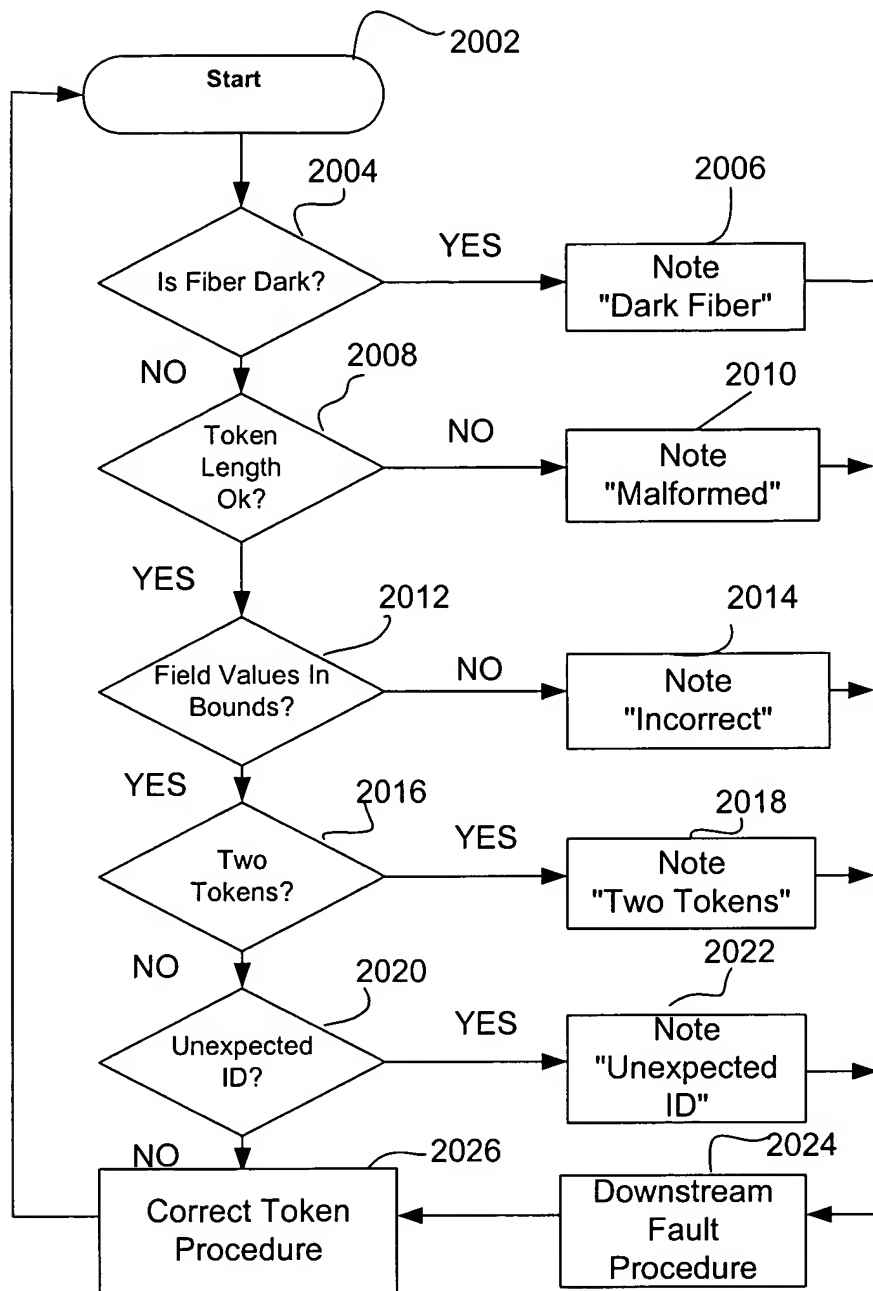


Figure 20

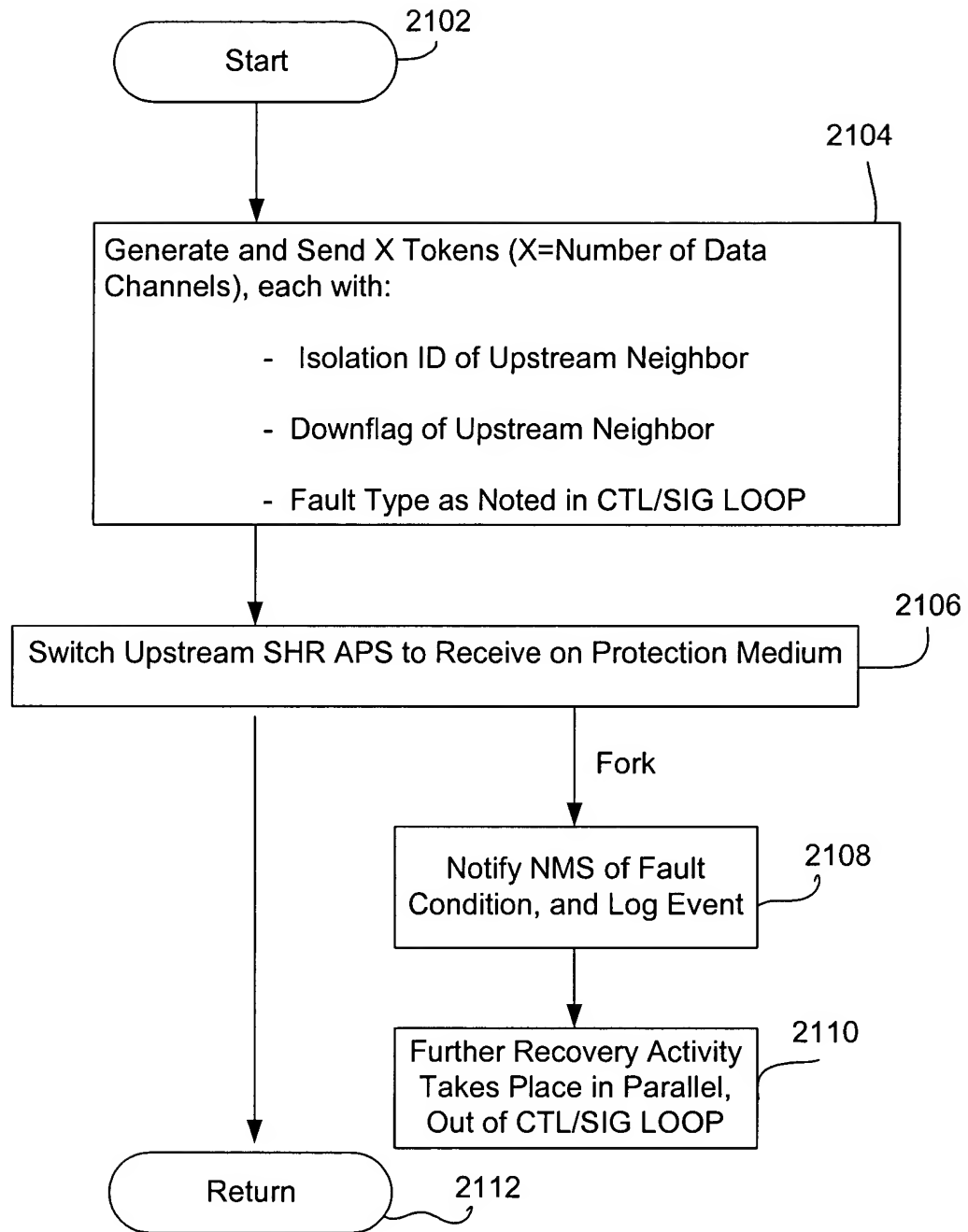


Figure 21

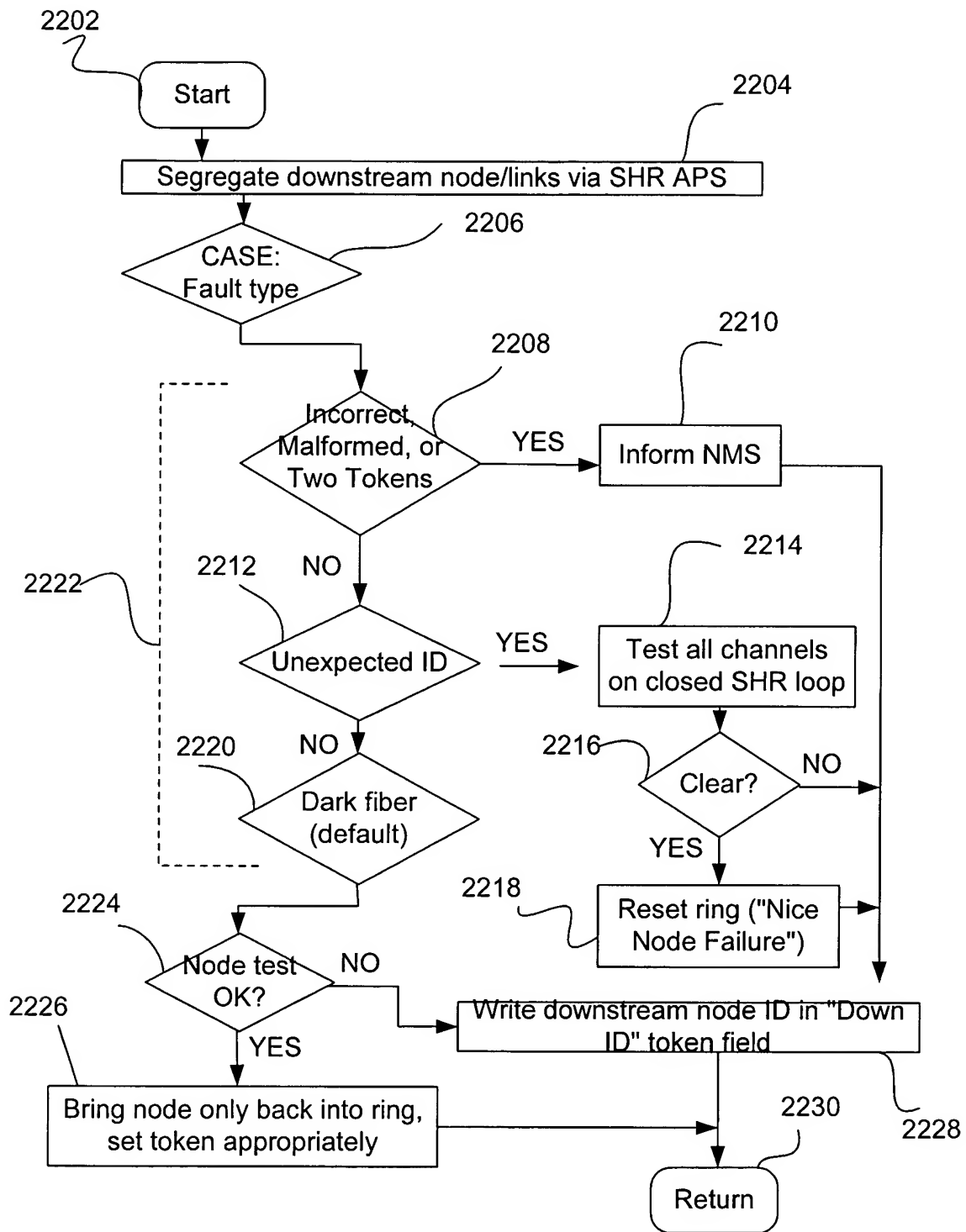


Figure 22

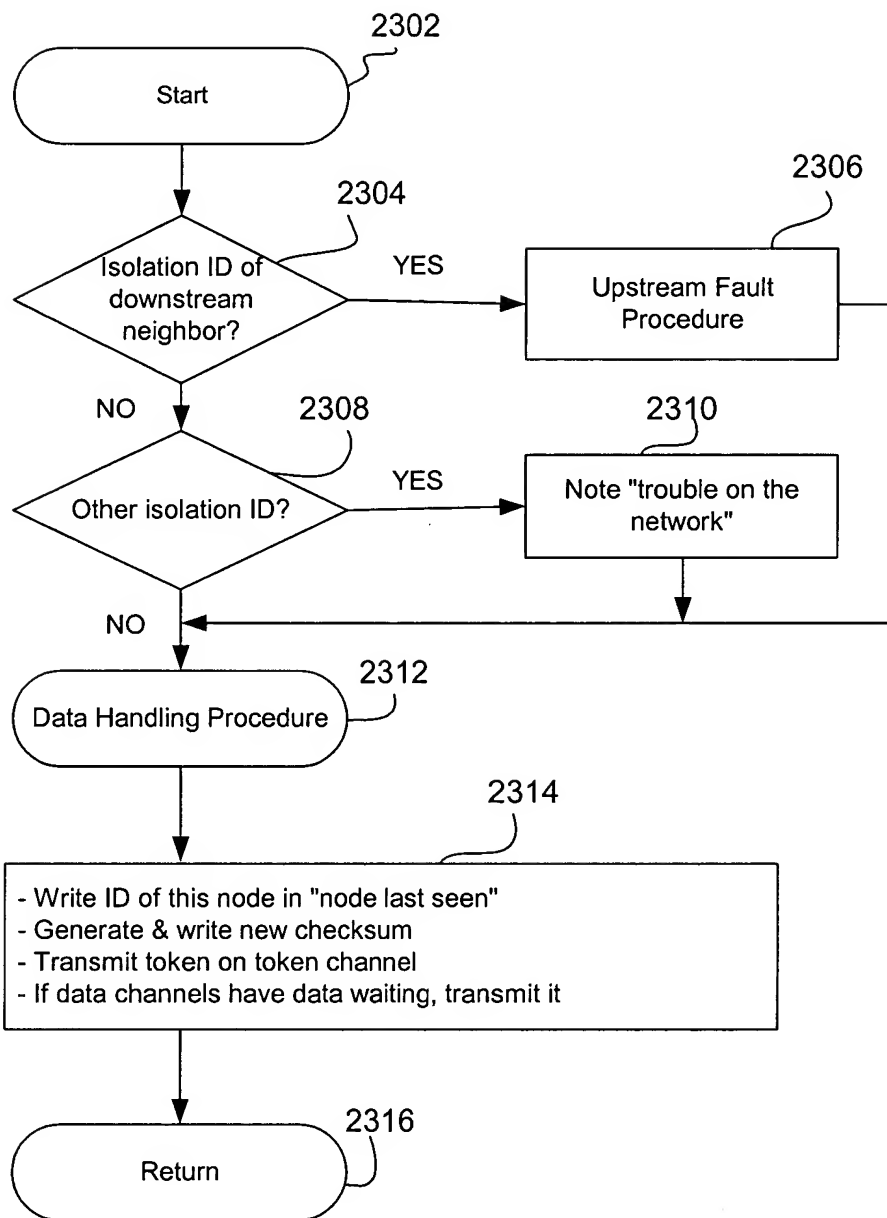


Figure 23

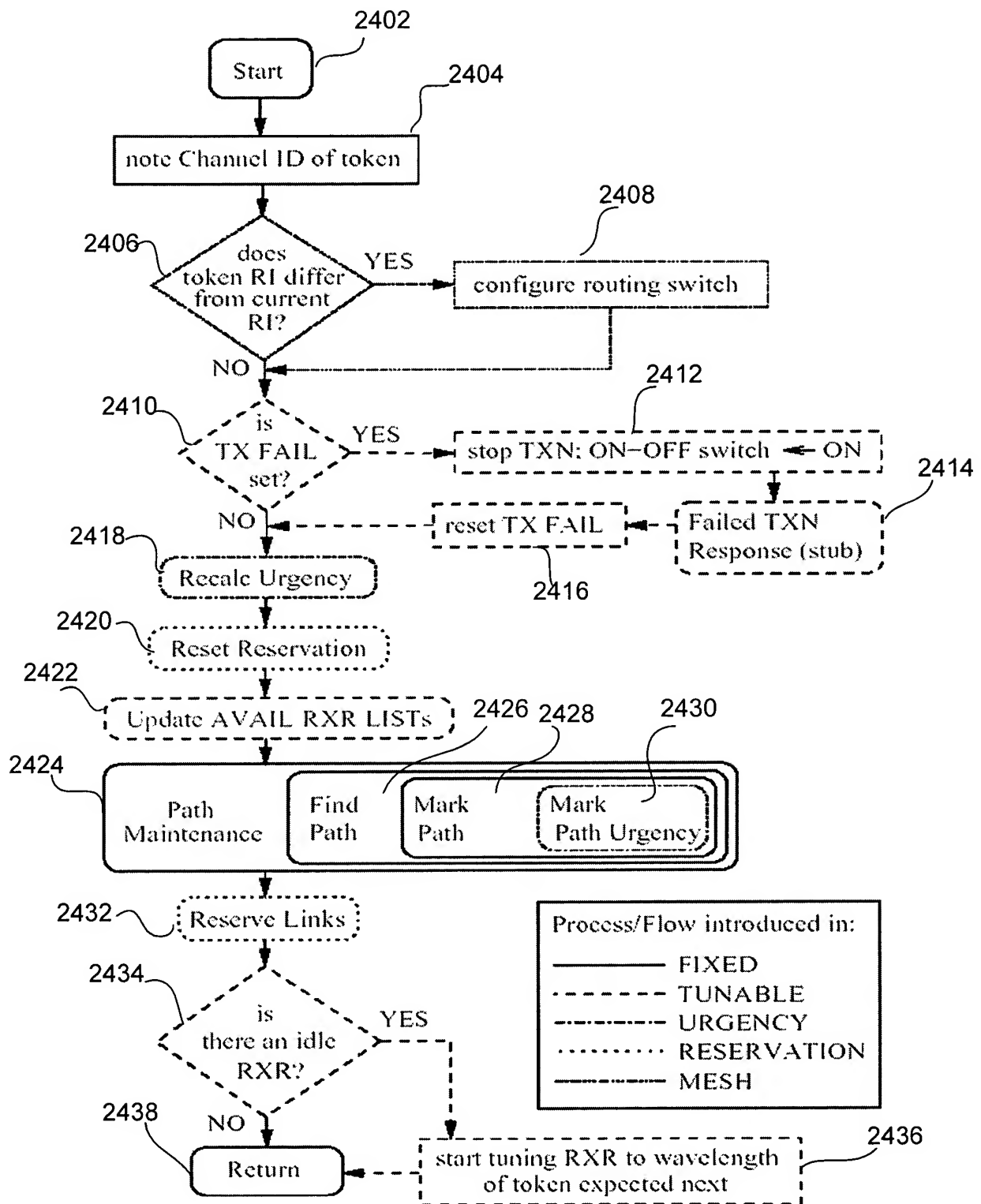
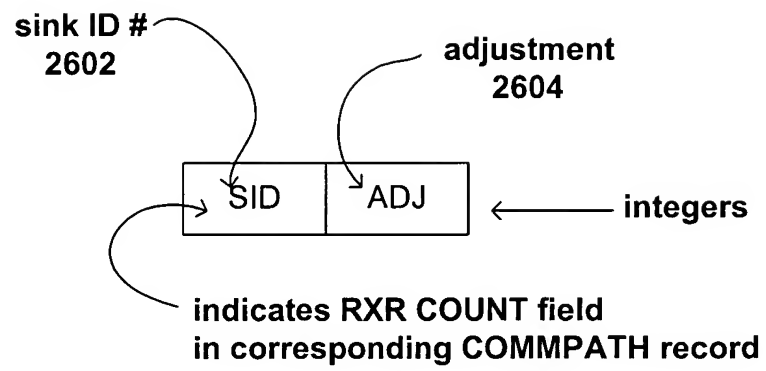


Figure 24



RXR COUNT LIST record fields

Figure 26

Algorithm 0.0.2: UPDATE AVAIL RXR LISTs(*Global, Node, Token*)

```

if rxr_lists are empty                                [block 0]
  then return

if Node is on a lightpath                               [block 1]
  then note source and sink

for each rec  $\in$  Node.add_back_rxr_list                 [block 2]
  {
    increment rec.adj
    increment Token[rec.sink].AVAIL_RXRS
    if Token[rec.sink].NUM_FAILS > 0
      then decrement Token[rec.sink].NUM_FAILS
    if rec.adj = 0
      then delete rec

for each rec1  $\in$  Node.take_away_rxr_list               [block 3]
  {
    decrement rec1.adj
    decrement Token[rec1.sink].AVAIL_RXRS
    if (Token[rec1.sink].AVAIL_RXRS + Token[rec1.sink].NUM_FAILS) < 0a
      {
        increment Token[rec1.sink].NUM_FAILS
        if sink noted and rec1.sink = sink and Token[rec1.sink].LINK = SINKb
          then if no active TXN to sink
            then { comment: TANDEM
                  Node.on_off[ $\lambda_i$ ]  $\leftarrow$  ON
                }
          else if Token[sink].LINK_URG  $\geq$  urgency of least urgent active TXN
            then { comment: STOMP
                  Node.on_off[ $\lambda_i$ ]  $\leftarrow$  ON
                  discontinue own least urgent active TXN
                  invoke Failed_TXN()
                }
          else { comment: SIPHON
                  reset lightpath from sink upstream
                  Node.on_off[ $\lambda_i$ ]  $\leftarrow$  OFF
                  Token[source].TX_FAIL  $\leftarrow$  sink
                }
        if rec1.adj = 0
          then { new rec2  $\leftarrow$  (rec1.sink, -(Global.num_tokens))
                add rec2 to Node.add_back_rxr_list
                delete rec1
              }
      }
  }

```

^aif the sum of the AVAIL_RXRS and NUM_FAILS token fields for *rec₁.sink* becomes negative ...

^bif *rec₁.sink* is the sink of a lightpath that was noted near the top of the algorithm ...

Figure 27

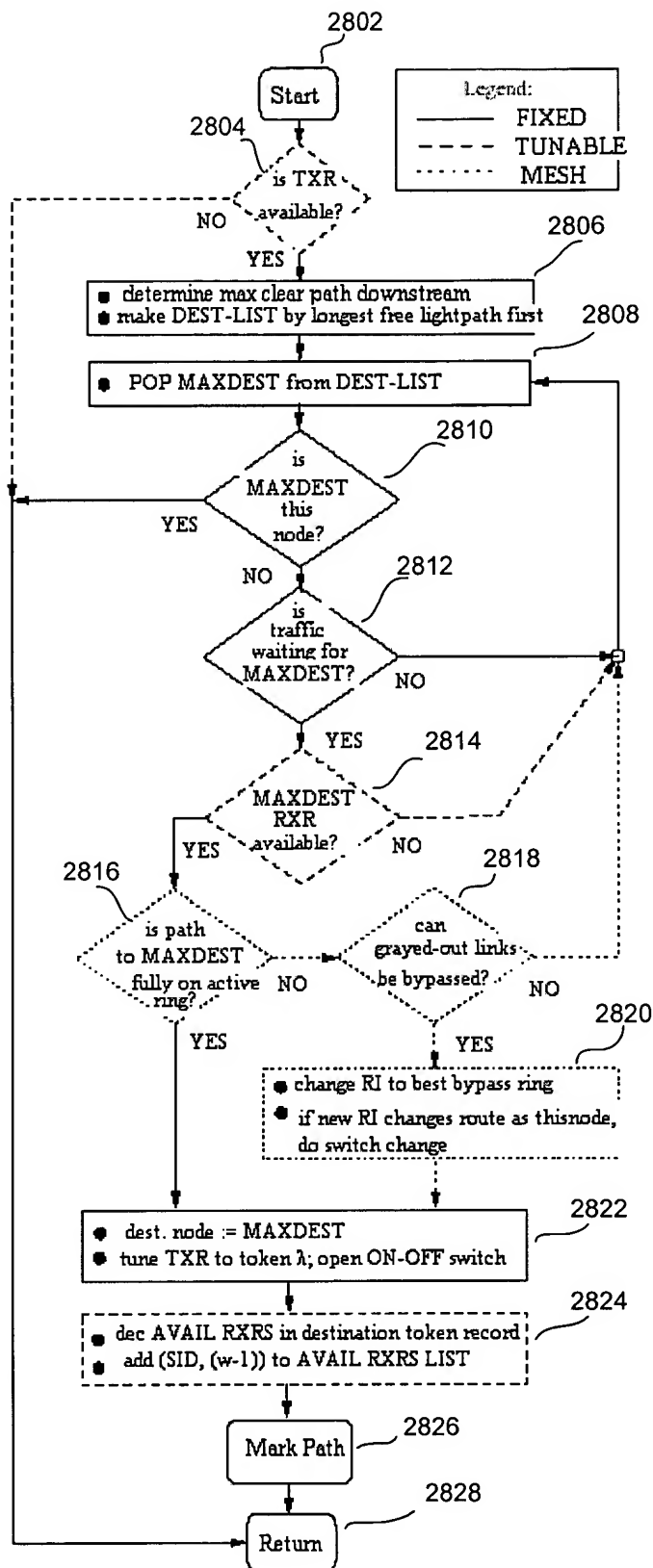


Figure 28

Algorithm 0.0.3: RESERVE LINKS PROCEDURE(*Token, Path*)

- create priority queue of destination *candidates*, sorted on primary key (most urgent),
and secondary key (greatest hop-length)

```
while (1)
    {
        if queue empty
            then return
    }
do {
    pop candidate
    if every RSV_URG of path to candidate has lower urgency than candidate burst
        then drop through WHILE loop
    }

for each link on Path
    {
        do {
            note “losing” reservation ID, if any
            write “my” ID and MAXDEST urgency number
        }

for each losing ID
        {
            do {
                reset all reservations which are contiguous to new reservation link
                and which have losing IDs (clear away “orphaned” IDs)
            }
        }
    }
```

Figure 29

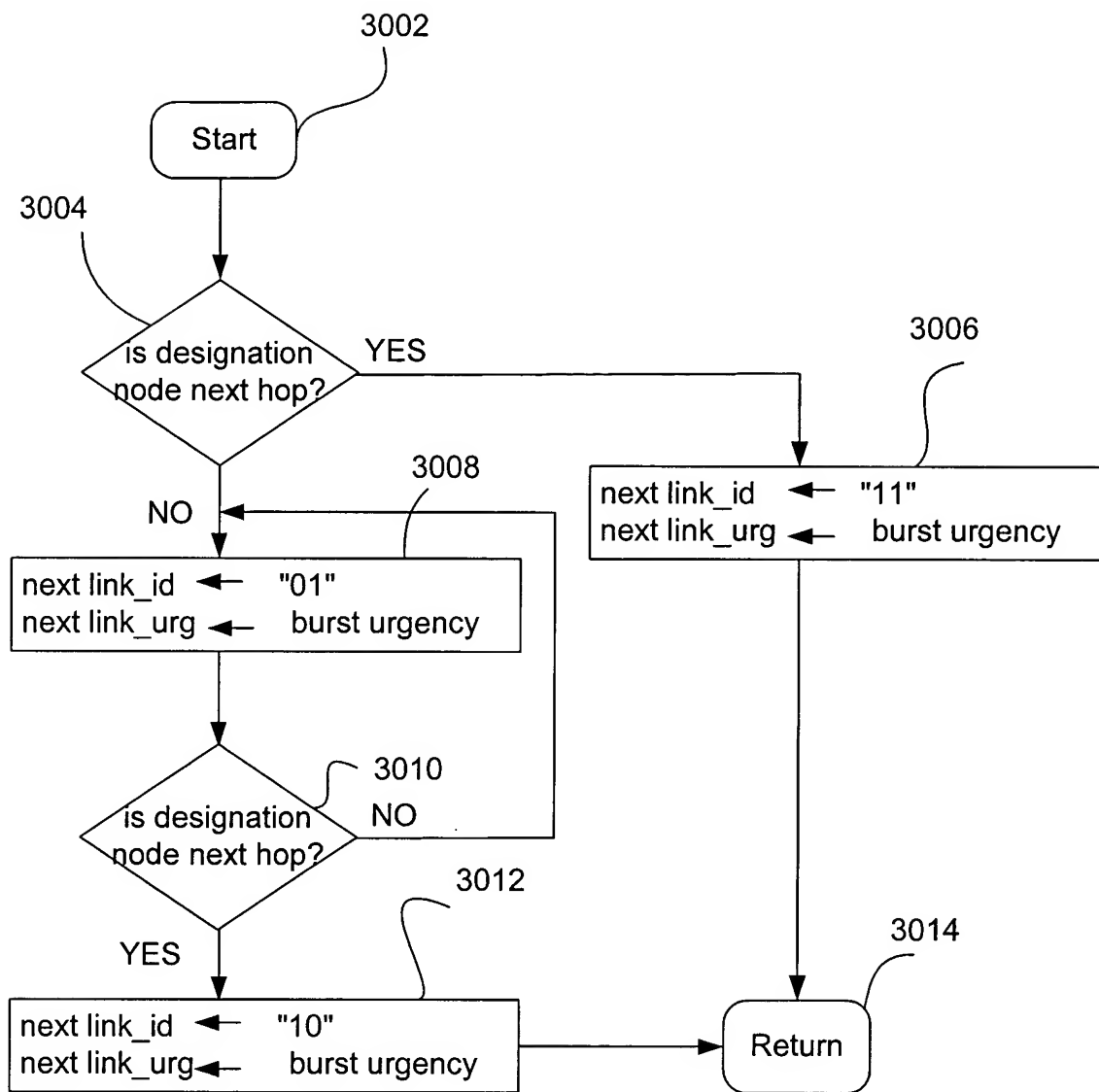


Figure 30

Algorithm 0.0.1: FIND PATH(*Global, Node, Token*)

if a TXR is available

```

    find max (the FREE link farthest downstream on active ring)
    dest_list ← all dests (with bursts waiting) incl. maxa
    sort dest_list, by primary key (most urgent)                                [3110]
                                     and secondary key (farthest)
    while (1)
        if dest_list is empty
            then return
        do {
            dest ← pop dest_list
            if (∀ intermediate link, (dest.urg > link.RSV_URG)                [3118]
                and “grayed-out” links can be bypassed)                        [3120]
                then breakb
        }
    if “grayed-out” links on path                                                [3130]
        then {
            Token.RI ← largest available bypass ring (RI)
            if new RI changes the route at thisnode, set switch
        }
    decrement Token[dest].AVAIL_RXRS
    arrec ← (dest, Global.num_tokens - 1)
    add arrec to Node.take_away_rxs_list
    mark_path(dest, dest.urg)

```

^aRecall that node information appears on the token in the same record with its upstream link.

^bThe break statement is unconditional, except in RESERVATION_SCHEME (first condition) and MESH (second condition).

Figure 31

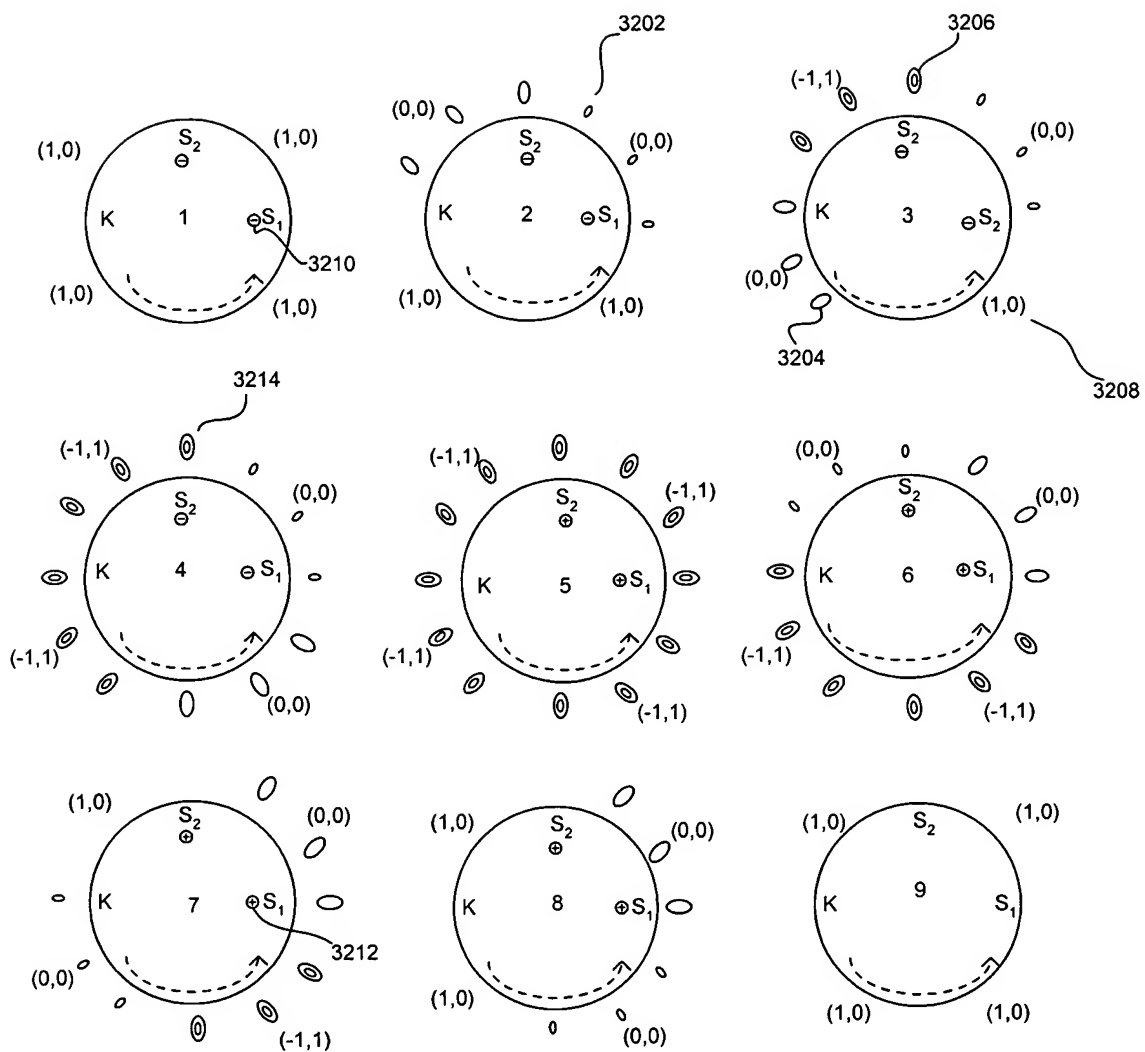
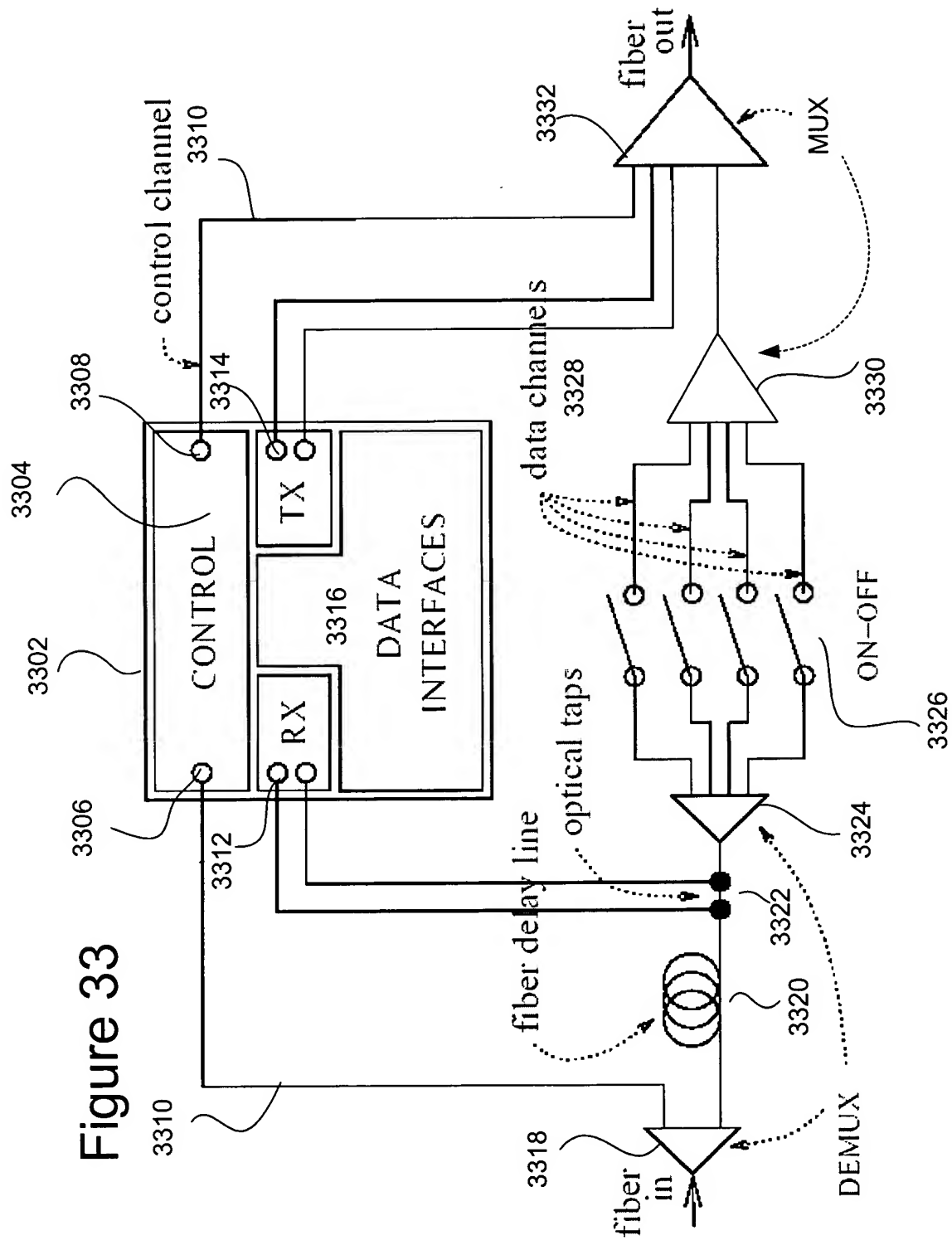


Figure 32

Figure 33

The diagram illustrates a fiber-optic communication system 3300. At the input, a fiber in 3318 enters a DEMUX 3324. The output of the DEMUX is a fiber delay line 3320, which is connected to an optical switch 3326. The optical switch 3326 consists of four ON-OFF switches. Optical taps 3322 are connected to the switch. The output of the optical switch is a fiber delay line 3312, which is connected to the RX block 3314 of a control unit 3302. The control unit 3302 also includes a CONTROL block 3304 and a TX block 3316. The TX block 3316 is connected to the MUX 3330 via a control channel 3310. The MUX 3330 is also connected to the fiber out 3332 via a control channel 3310.



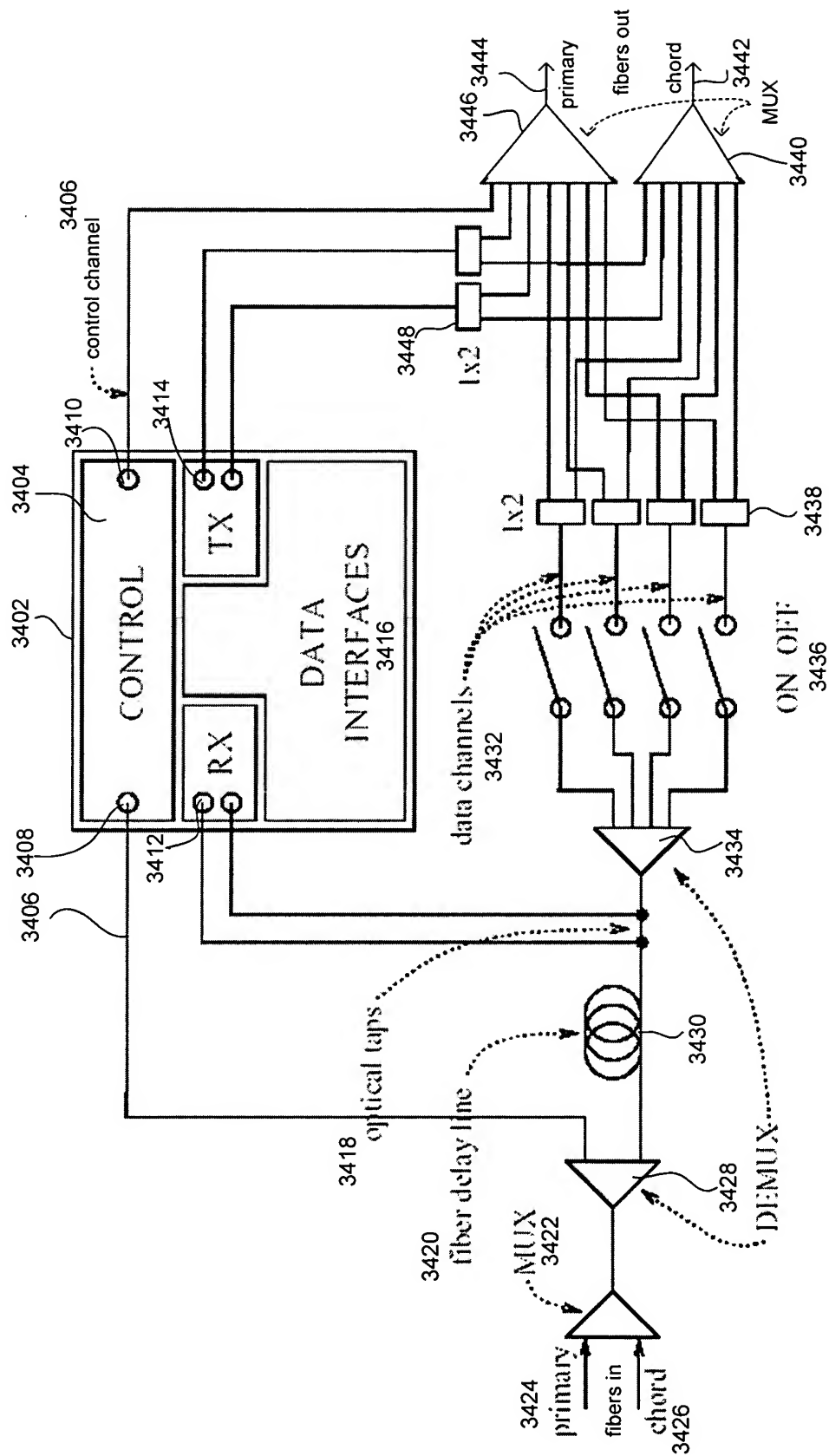


Figure 34